## Review of Lecture 5

- Dichotomies

- Growth function

$$
m_{\mathcal{H}}(N)=\max _{\mathbf{x}_{1}, \cdots, \mathbf{x}_{N} \in \mathcal{X}}\left|\mathcal{H}\left(\mathbf{x}_{1}, \cdots, \mathbf{x}_{N}\right)\right|
$$

- Break point

- Maximum \# of dichotomies

| $\mathbf{X}_{1}$ | $\mathbf{X}_{2}$ | $\mathbf{X}_{3}$ |
| :---: | :---: | :---: |
| $\circ$ | $\circ$ | $\circ$ |
| $\circ$ | $\circ$ | $\bullet$ |
| $\circ$ | $\bullet$ | $\circ$ |
| $\circ$ | $\circ$ | $\circ$ |

# Learning From Data 

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## Lecture 6: Theory of Generalization

## Outline

- Proof that $m_{\mathcal{H}}(N)$ is polynomial
- Proof that $m_{\mathcal{H}}(N)$ can replace $M$

Bounding $m_{\mathcal{H}}(N)$

To show: $\quad m_{\mathcal{H}}(N)$ is polynomial

We show: $\quad m_{\mathcal{H}}(N) \leq \cdots \leq \cdots \leq$ a polynomial

Key quantity:
$B(N, k)$ : Maximum number of dichotomies on $N$ points, with break point $k$

## Recursive bound on $B(N, k)$

Consider the following table:
$B(N, k)=\alpha+2 \beta$

|  | \# of rows | $\mathrm{x}_{1}$ $\mathrm{x}_{2}$ | $\ldots \mathrm{x}_{N-1}$ | $\mathrm{x}_{N}$ |
| :---: | :---: | :---: | :---: | :---: |
| $S_{1}$ | $\alpha$ | +1 +1 | .. +1 | +1 |
|  |  | $-1+1$ | +1 | -1 |
|  |  | : |  | : |
|  |  | +1 $\quad-1$ | ... -1 | -1 |
|  |  | -1 +1 | ... -1 | +1 |
| $S_{2} S_{2}^{+}$ | $\beta$ | +1 -1 | $\ldots \quad+1$ | $+1$ |
|  |  | $\begin{array}{ll}-1 & -1\end{array}$ | +1 | +1 |
|  |  | : | : : | : |
|  |  | +1 -1 | $\ldots$. +1 | +1 |
|  |  | -1 -1 | $\ldots$... 1 | +1 |
| $S_{2}^{-}$ | $\beta$ | +1 $\quad-1$ | $\ldots \quad+1$ | -1 |
|  |  | $\begin{array}{ll}-1 & -1\end{array}$ | $\ldots$. +1 | -1 |
|  |  | - | : | : |
|  |  | +1 -1 | $\ldots \quad+1$ | -1 |
|  |  | -1 -1 | $\ldots$.. -1 | -1 |

## Estimating $\alpha$ and $\beta$

Focus on $\mathbf{x}_{1}, \mathbf{x}_{2}, \cdots, \mathbf{x}_{N-1}$ columns:

$$
\alpha+\beta \leq B(N-1, k)
$$



## Estimating $\beta$ by itself

Now, focus on the $S_{2}=S_{2}^{+} \cup S_{2}^{-}$rows:

$$
\beta \leq B(N-1, k-1)
$$



## Putting it together

$$
B(N, k)=\alpha+2 \beta
$$

$$
\begin{aligned}
& \alpha+\beta \leq B(N-1, k) \\
& \beta \leq B(N-1, k-1)
\end{aligned}
$$

$$
B(N, k) \leq
$$

$$
B(N-1, k)+B(N-1, k-1)
$$

|  | \# of rows | $\mathrm{x}_{1}$ $\mathrm{x}_{2}$ | $\ldots \mathrm{x}_{N-1}$ | $\mathrm{x}_{N}$ |
| :---: | :---: | :---: | :---: | :---: |
| $S_{1}$ | $\alpha$ | +1 +1 | +1 | +1 |
|  |  | $-1+1$ | $\ldots$ +1 | -1 |
|  |  | : : | : | : |
|  |  | +1 -1 | ... -1 | -1 |
|  |  | -1 +1 | ... -1 | +1 |
| $S_{2}^{+}$ | $\beta$ | +1 -1 | $\ldots{ }^{\text {. }}$ +1 | +1 |
|  |  | $\begin{array}{ll}-1 & -1\end{array}$ | +1 | +1 |
|  |  | : | : : | : |
|  |  | +1 -1 | ... +1 | +1 |
|  |  | -1 -1 | $\ldots-1$ | +1 |
| $S_{2}$ | $\beta$ | +1 -1 | .. +1 | -1 |
|  |  | $\begin{array}{ll}-1 & -1\end{array}$ | $\ldots$ +1 | -1 |
|  |  | : | : | : |
|  |  | +1 -1 | ... +1 | -1 |
|  |  | -1 -1 | .. -1 | -1 |

Numerical computation of $B(N, k)$ bound

$$
B(N, k) \leq B(N-1, k)+B(N-1, k-1)
$$

|  |  | $k$ |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 1 | 2 | 3 | 4 | 5 | 6 | $\ldots$ |  |
|  | 1 | 1 | 2 | 2 | 2 | 2 | 2 |  |$]$.

Analytic solution for $B(N, k)$ bound
$B(N, k) \leq B(N-1, k)+B(N-1, k-1)$

Theorem:

$$
B(N, k) \leq \sum_{i=0}^{k-1}\binom{N}{i}
$$

1. Boundary conditions: easy

|  |  | $k$ |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | 1 | 2 | 3 | 4 | 5 | 6 | $\ldots$ |
|  | 1 | 1 | 2 | 2 | 2 | 2 | 2 | $\cdots$ |
|  | 2 | 1 |  |  |  |  |  |  |
|  | 3 | 1 |  |  |  |  |  |  |
| $N$ | 4 | 1 |  |  |  |  | $\bullet$ |  |
|  | 5 | 1 |  |  |  |  |  |  |
|  | 6 | 1 |  |  |  |  |  |  |
|  | $:$ | $:$ |  |  |  |  |  |  |
|  |  |  |  |  |  |  |  |  |

## 2. The induction step

$$
\begin{aligned}
\sum_{i=0}^{k-1}\binom{N}{i} & =\sum_{i=0}^{k-1}\binom{N-1}{i}+\sum_{i=0}^{k-2}\binom{N-1}{i} ? \\
& =1+\sum_{i=1}^{k-1}\binom{N-1}{i}+\sum_{i=1}^{k-1}\binom{N-1}{i-1} \\
& =1+\sum_{i=1}^{k-1}\left[\binom{N-1}{i}+\binom{N-1}{i-1}\right] \\
& =1+\sum_{i=1}^{k-1}\binom{N}{i}=\sum_{i=0}^{k-1}\binom{N}{i} \checkmark
\end{aligned}
$$

## It is polynomial!

For a given $\mathcal{H}$, the break point $k$ is fixed

$$
m_{\mathcal{H}}(N) \leq \underbrace{\sum_{i=0}^{k-1}\binom{N}{i}}_{\text {maximum power is } N^{k-1}}
$$

Three examples

$$
\sum_{i=0}^{k-1}\binom{N}{i}
$$

- $\mathcal{H}$ is positive rays: (break point $k=2$ )

$$
m_{\mathcal{H}}(N)=N+1 \leq N+1
$$

- $\mathcal{H}$ is positive intervals: (break point $k=3$ )

$$
m_{\mathcal{H}}(N)=\frac{1}{2} N^{2}+\frac{1}{2} N+1 \leq \frac{1}{2} N^{2}+\frac{1}{2} N+1
$$

- $\mathcal{H}$ is 2 D perceptrons: (break point $k=4$ )

$$
m_{\mathcal{H}}(N)=? \leq \frac{1}{6} N^{3}+\frac{5}{6} N+1
$$

## Outline

- Proof that $m_{\mathcal{H}}(N)$ is polynomial
- Proof that $m_{\mathcal{H}}(N)$ can replace $M$


## What we want

## Instead of:

$$
\mathbb{P}\left[\left|E_{\text {in }}(g)-E_{\text {out }}(g)\right|>\epsilon\right] \leq 2 \quad M \quad e^{-2 \epsilon^{2} N}
$$

We want:

$$
\mathbb{P}\left[\left|E_{\text {in }}(g)-E_{\text {out }}(g)\right|>\epsilon\right] \leq 2 m_{\mathcal{H}}(N) e^{-2 \epsilon^{2} N}
$$

## Pictorial proof

- How does $m_{\mathcal{H}}(N)$ relate to overlaps?
- What to do about $E_{\text {out }}$ ?
- Putting it together


What to do about $E_{\text {out }}$

-ゃ००००००००
$E_{\text {in }}(h)$

$E_{\text {in }}(h) \bullet \bullet \bullet \bullet \bullet \bullet \bullet \bullet \bullet \bullet ~$ $E_{\text {in }}^{\prime}(h) \bullet \bullet \bullet \bullet \bullet \bullet \bullet \bullet \bullet \bullet$

## Putting it together

Not quite:

$$
\mathbb{P}\left[\left|E_{\text {in }}(g)-E_{\text {out }}(g)\right|>\epsilon\right] \leq 2 m_{\mathcal{H}}(N) e^{-2 \epsilon^{2} N}
$$

but rather:

$$
\mathbb{P}\left[\left|E_{\text {in }}(g)-E_{\text {out }}(g)\right|>\epsilon\right] \leq 4 m_{\mathcal{H}}(2 N) e^{-\frac{1}{8} \epsilon^{2} N}
$$

The Vapnik-Chervonenkis Inequality

